

## The Leverhulme Research Centre for Functional Materials Design

# Crystal Structure Prediction by Vertex Removal in Euclidean Space

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# What is Crystal Structure Prediction?

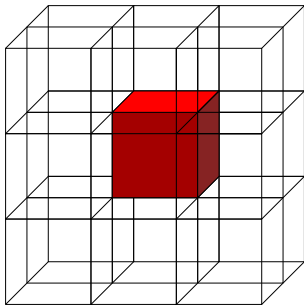
**Problem** *Crystal Structure Prediction (CSP)*

**Input:** A set of ions,  $A$ , an area of space,  $C$ .

**Output:** A structure,  $S$ , made by placing some copies of the ions in  $A$  in  $C$ , with a *neutral* charge *minimising* potential energy between the ions.

## What is a *Crystal*?

- We consider crystals to be made up of unit cells.
- Each unit cell is the smallest repeating region of space within the crystal.



**Figure 1:** Unit cell highlighted in red, note any other box would be equivalent

## What is a *Unit Cell*?

- Each unit cell is a collection of *ions*.
- We assume each unit cell is *independent* of all other unit cells.
  - This means that we only consider the *interaction* of ions within the same cell.
- Every cell must have a total *Neutral* charge.

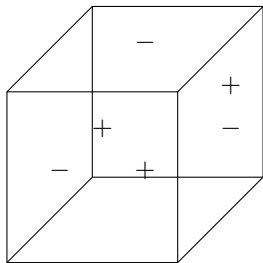


Figure 2: Example of a unit cell

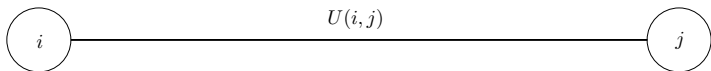
## What is an *Ion*?

- A point in space belonging to a *species*.
- The species determines its interaction with other ions, as well as its charge.
- We denote the charge of ion,  $i$ ,  $q_i$
- for a unit cell with the set of ions,  $S$ , we require the following to be satisfied:

$$\sum_{i \in S} q_i = 0$$

## How do we determine *interaction*

- We define the pairwise interaction for any pair of ions by some parameterised function  $U_{\theta_{i,j}}(r_{ij})$ .
- $r_{ij}$  is the distance between ions  $i$  and  $j$
- The parameters for this function are determined by the species of ions  $i$  and  $j$ .
- A negative value for interaction means that the ions are trying to move closer together, which implies the crystal will be stronger.
- We can use this to represent the ions as a graph with weighted edges.

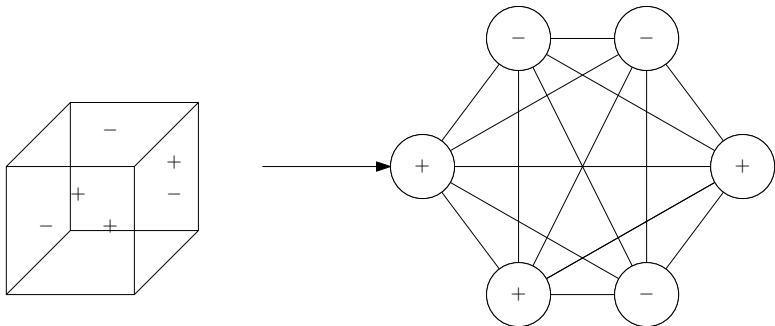


**Figure 3:** Interaction between two ions represented as a graph, each ion represents a vertex.

## Representing the unit cell as a graph

- We can use this to represent the unit cell as a graph embedded into 3-dimensional space.
- Conversely, we can use this to create a complete graph where:
  - each ion is a vertex.
  - each edge has a weight equal to the interaction between the two ions it is connected to.

## Representing the unit cell as a graph



**Figure 4:** We can redraw our unit cell as a graph, using the ions as vertices and the interactions as weights on the edges.



## Some Notation and Terminology

- A *Structure* refers to a set containing all the ions within a given unit cell.
- Give a structure,  $S$ , we use  $S^+$  to denote the set of ions with a positive charge, and  $S^-$  for the ions with a negative charge.
- We use  $|S^+|$  to denote the sum of the magnitude of the positive charges,  $|S^+| = \sum_{i \in S^+} q_i$
- We use  $|S^-|$  to denote the sum of the magnitude of the negative charges,  $|S^-| = \sum_{i \in S^-} -(q_i)$ .

## Crystal Structure Prediction by Vertex Removal

- We can use this to define our problem: Crystal Structure Prediction by K-Ion (Vertex) Removal.
- We take as input some highly dense initial structure,  $S = S^+ \cup S^-$ , within our unit cell, and an integer,  $k$ .
- Our goal is to remove some substructure of  $S$ ,  $S = -S'^+ \cup S'^-$ , such that:

$$|S'^+| \leq |S^+| - k$$

$$|S'^-| \leq |S^-| - k$$

$$|S'^+| = |S'^-|$$

- We also want our solution to be *minimal*, in that there is no substructure,  $S'' \subset S'$ , that also satisfies the above.

# Crystal Structure Prediction by Vertex Removal

**Problem** *K-Vertex Removal* (KIR).

*Input* A structure of ions,  $S$ , a pairwise energy function,  $U$ , and an integer  $k$ .

*Output* A substructure,  $S' \subseteq S$ , formed by a minimal removal of  $k$  charges from  $S$  with minimal total energy with respect to  $U$

## Decision problem

- For NP-Completeness, we need to reformulate this as a decision problem.
- We do this by adding a goal energy,  $g$ , which is the maximum allowed energy.
- Given an instance of  $KIR$ , we report yes if there is a substructure with total energy less than or equal to  $g$ , or no otherwise.

## The Energy function

- The given energy function determines what will and won't be a good solution.
- We will be considering a general class of functions for which this problem is NP and APX complete, which we call the *Crystalline* class of functions,  $\mathcal{F}$ .

$\forall f \in \mathcal{F}, \exists a, b \in \mathbb{R}, a > b$  s.t.

$$\forall r \in \mathbb{R}^+ \exists \theta_{\mathbf{ar}}, \theta_{\mathbf{br}} \in \mathbb{R}^n \text{ s.t. } f_{\theta_{\mathbf{ar}}}(r) = a, f_{\theta_{\mathbf{br}}} = b$$

## Buckingham-Coulomb potential energy

- The *Buckingham-Coulomb* potential is used frequently in computational chemistry for determining the energy between ions.
- In this function we use the charge of the ions, as well as 3 *force field* parameters, determined by the species of the ions, as our parameters.
- These are  $A_{ij}$ ,  $B_{ij}$  and  $C_{ij}$ .
- The energy function is:

$$U_{\{A_{ij}, B_{ij}, C_{ij}, q_i, q_j\}}(r_{ij}) = \frac{A_{ij}}{e^{B_{ij} r_{ij}}} - \frac{C_{ij}}{r_{ij}^6} + \frac{q_i q_j}{r_{ij}}$$

## Clique to K-Vertex removal, for $U \in \mathcal{F}$

- Assume we have some instance of the Clique problem, where we have a graph,  $G$ , and wish to find a clique of size  $k$ .
- We claim that we can reduce this problem to KIR, making the latter NP and APX complete.
- We will do this by constructing a structure such that we will be left with only the ions corresponding to vertices in  $G$  in a clique of size  $k$ .
- The main idea is to create 2 ions for each vertex in  $G$ , labelled with their corresponding vertex.
- We will assume that our energy function is some arbitrary function in  $\mathcal{F}$

# Parameters

- We assign parameters so that the energy between a pair of ions,  $i, j$  corresponding to  $v_i$  and  $v_j$ , is as follows:

$$U_{\theta_{ij}}(r_{ij}) = \begin{cases} b & \text{if } (v_i, v_j) \in E, \text{ or } v_i = v_j \\ a & \text{otherwise} \end{cases}$$

- We know from our definition of  $\mathcal{F}$  that we can always achieve this.