Principles of Computer Game Design and Implementation

Lecture 27

Outline for today

• pathfinding

The Problem

Pathfinding

- Given the current position and the target position
 - Calculate a sequence of positions (path)
 - Can follow with steering
 - Shortest / lowest cost path



Pathfinding In Games

- Initially the concept was only used in RTS
- Now *the* most important AI technique
 (probably)
- Still can be buggy. See
 - <u>http://www.ai-blog.net/archives/000152.html</u>
 - <u>http://www.youtube.com/watch?v=lw9G-</u>
 <u>8gL5o0&feature=player_embedded</u>

Tackling Paths

- Characters "live" in a computer world
 - Even developers may not know exact location
 - Physics simulations





• Pathfinders operate on discrete structures

Remember This?



Romania Map



From COMP219:

- A *search* algorithm can solve the navigation problem
- Simple algorithms
 - Breadth-first, depth-first, unit cost,...
 do not work in real-world problems
- A* is the best we have

So

- A* works on *weighted graphs*
 - Pathfinding graphs
 - Explicitly or implicitly represented
 - Romania map: explicit representation
 - Many games do not store full graphs



Recall: Search Tree



- An imaginary tree showing all possible states reachable from the initial state
- Search *strategy* defines an expansion order

Recall: A* Search (Strategy)

- Combine uniform cost search and greedy search.
- Uses *heuristic* f:

f(n) = g(n) + h(n),

- where
 - g(n) is path cost of n;
 - h(n) is expected cost of cheapest solution from n.

Recall: General Algorithm for A* Search

agenda = initial state; while agenda not empty do take node from agenda such that $f(node) = min \{ f(n) | n in agenda \}$ where f(n) = q(n) + h(n)if node is goal state then return solution; new nodes = apply operations to node; add new nodes to the agenda;



The general framework allows to visit nodes more than once

• **Closed** nodes list: already visited nodes

Theory V Practice: Admissible and Inadmissible Heuristics

- A* is guided by heuristic
- If heuristic is too high (overestimates)
 - It's **inadmissible**
 - A* is not guaranteed to find best path
 - Does not mean you cannot use it!
 - Faster search vs better paths balance
 - Closed nodes can be "reopened"

A* Requires

- To store the agenda
 Open nodes list
- To store the
 Closed nodes list

- For every open node: costs so far and estimated costs
- For every closed node the *connection* (edge) leading to it

Pathfinding Algorithm

while lowest rank in open is not goal
 current = remove lowest rank item from open;
 closed.add(current);

for neighbors of current:

Ncost = g(current) + cost(current, neighbor);

- if (open.contains(neighbor)&&Ncost<g(neighbor))
 open.remove(neighbor)</pre>
- if (closed.contains(neighbor)&&Ncost<g(neighbor))
 closed.remove(neighbor)</pre>
- if (!open.contains(neighbor) & &

!closed.contains(neighbor))

```
g(neighbor) = Ncost
```

open.add(neighbour)

neighbor.connection = current

Good Practice: Class GraphNode

public class GraphNode { // link to game world Vector<Edge> edges } public class Edge { GraphNode from, to; float cost;

Good Practice: NodeRecord

public class NodeRecord {

GraphNode node;

Edge connection;

float costSoFar;

float estimatedGoalCost;

float currentCost;

}

Data Structures

- **Closed**: unsorted list of NodeRecord
- Open
 - Unsorted list of NodeRecord
 - Insert: easy (just append)
 - Take: hard (loop through all of them)
 - Priority queue of NodeRecord
 - Insert: medium (balancing)
 - Take: medium



Simplicity Rules

- On a grid-like graph
 One take per 8 inserts
- With a good heuristics
 - A simple unsorted list might be more efficient than a sophisticated Priority Queue!



Tile-Based Games

- A vast majority of RTS games are tile-based
 - Every unit occupies (one or more) tile
 - Every tile can accommodate \leq 1 unit
- A tile is either **blocked** or **passable**



Tile Shapes

• Different games use different tiles



Nodes

- A node is uniquely identified with (x,y) coordinates
- No need to store neighbour nodes

Easily compute when necessary

Heuristics



• Manhattan block distance: $\Delta x + \Delta y$



Trouble: too many paths of same value

Breaking Ties

- Breaking ties is one of the reasons to consider an *inadmissible heuristics:*
 - Biased towards pursuing the goal
 - A* can run faster
 - If it is just slightly higher, A* will still find best paths
- Other reason?
 - Distance in hours, heuristics in km
 - Computational complexity

Heuristics



• Diagonal moves allowed: $\Delta x + \Delta y$



Heuristics



• Euclidian distance: $\sqrt{(\Delta x)^2 + (\Delta y)^2}$



Worst Possible Case

- Worst possible case for any search algorithm
 - No path



- Will explore all available space

Updated Pathfinding

- Check if Start and Finish are valid locations
 - If Finish is not valid, no path
 - If **Start** is not valid
 - Something goes wrong
 - Delete agent?
 - Move to a valid location?
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Zone Mapping

- Every tile belongs to a zone
 - 0 impassable
 - Same number –

can pass



- Zone equivalence array

Hovercraft[0]=0; Hovercraft[1]=0; Hovercraft[2]=0

Zone Equivalence Array

- For every zone number and
- Every vehicle class
 - ZEA[zone number]
 - Either zone itself
 - Or the smallest equivalent zone number
- If (ZEA[S.zone] == ZEA[F.zone])
 Call the pathfinder

Pathfinding Pool

- Running an A* algorithm takes time
- In RTS games there are dozens of characters
- If every one of them starts A*...
 - A *pool* of pathfinders
 - A queue of agents waiting for paths
 - Start moving / play animation while waiting